

# Funkce vektorové náhodné proměnné

náhodná proměnná  $\mathbf{x} = (x_1, x_2, \dots, x_m)$  popsána celkovou hustotou pravděpodobnosti  $f(\mathbf{x})$

nová náhodná proměnná  $\mathbf{a} = (a_1, a_2, \dots, a_m)$  vytvořená transformací  $\mathbf{a} = \mathbf{h}(\mathbf{x})$

$$\begin{array}{llll} a_1 & = & h_1(x_1, x_2, \dots, x_m) & \text{zpětná transformace} & x_1 & = & h_1^{-1}(a_1, a_2, \dots, a_m) \\ \dots & & & & \dots & & \\ a_m & = & h_m(x_1, x_2, \dots, x_m) & & x_m & = & h_m^{-1}(a_1, a_2, \dots, a_m) \end{array}$$
$$\mathbf{x} = \mathbf{h}^{-1}(\mathbf{a})$$

celková hustota pravděpodobnosti nové náhodné proměnné  $\mathbf{a}$  je  $g(\mathbf{a}) = f(\mathbf{h}^{-1}(\mathbf{a})) |\mathbf{J}|$ ,

kde  $\mathbf{J}$  je Jacobiho matice

$$\mathbf{J} = \begin{pmatrix} \frac{\partial h_1^{-1}}{\partial a_1} & \frac{\partial h_1^{-1}}{\partial a_2} & \dots & \frac{\partial h_1^{-1}}{\partial a_m} \\ \dots & \dots & \dots & \dots \\ \frac{\partial h_m^{-1}}{\partial a_1} & \frac{\partial h_m^{-1}}{\partial a_2} & \dots & \frac{\partial h_m^{-1}}{\partial a_m} \end{pmatrix}$$

# Funkce náhodné proměnné – konvoluce

náhodné proměnné  $(x, y)$  s celkovou hustotou pravděpodobnosti  $f(x, y)$

nová náhodná proměnná  $u = x + y$

pomocná proměnná  $v = y$

$$u = x + y$$

zpětná transformace

$$x = u - v$$

$$v = y$$

$$y = v$$

celková hustota pravděpodobnosti nových náhodných proměnných  $g(u, v) = f(u - v, v)|\mathbf{J}|$

$$\mathbf{J} = \begin{pmatrix} \frac{\partial x}{\partial u} & \frac{\partial x}{\partial v} \\ \frac{\partial y}{\partial u} & \frac{\partial y}{\partial v} \end{pmatrix} = \begin{pmatrix} 1 & -1 \\ 0 & 1 \end{pmatrix}$$

$$|\mathbf{J}| = 1 \longrightarrow g(u, v) = f(u - v, v)$$

marginální hustota pravděpodobnosti

$$g_u(u) = \int_{-\infty}^{\infty} g(u, v) dv = \int_{-\infty}^{\infty} f(u - v, v) dv$$

pokud jsou náhodné proměnné  $x, y$  nezávislé

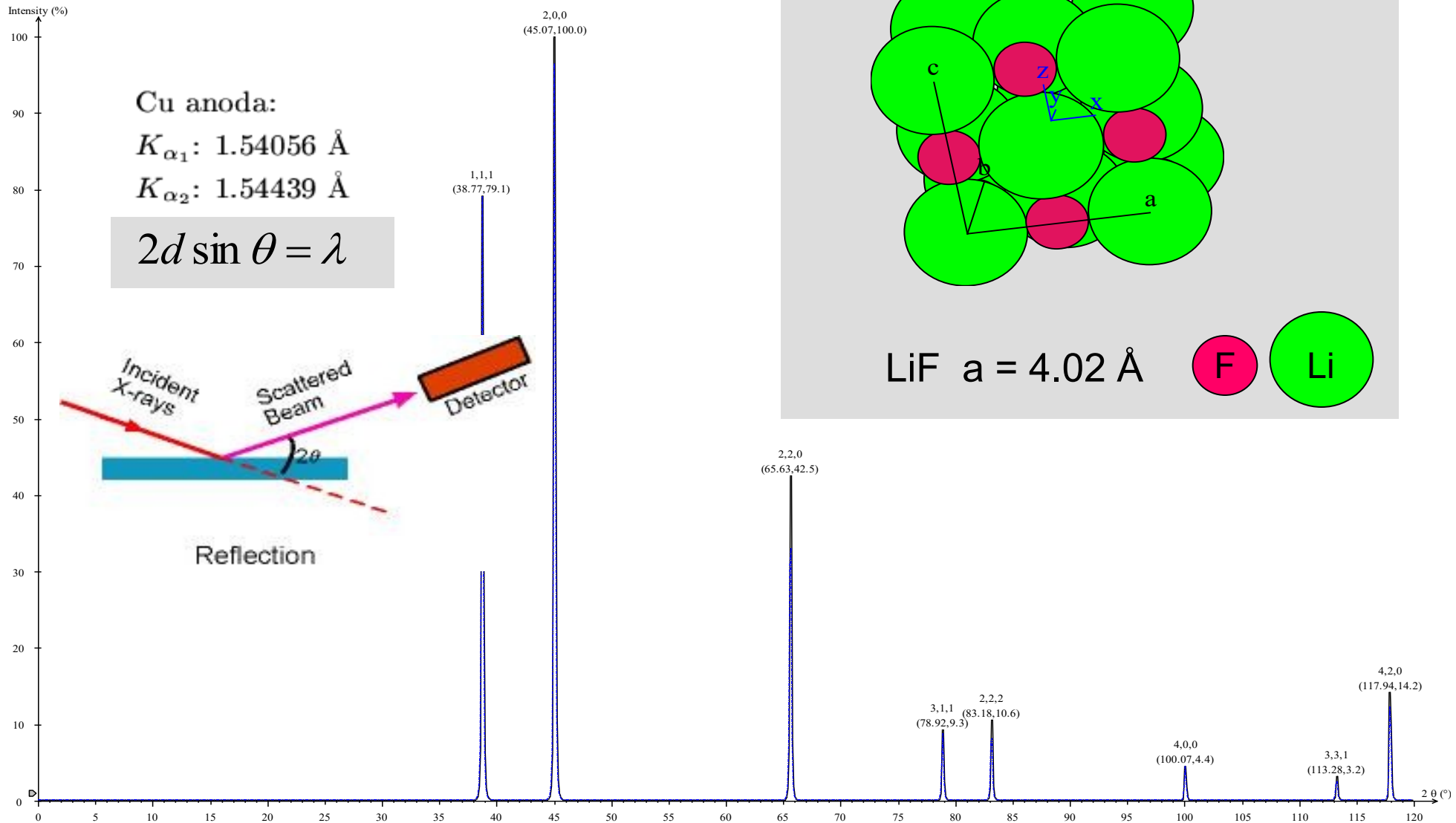
$$f(x, y) = f_x(x)f_y(y)$$

$$g_u(u) = \int_{-\infty}^{\infty} f_x(u - v)f_y(v) dv = f_x \star f_y$$

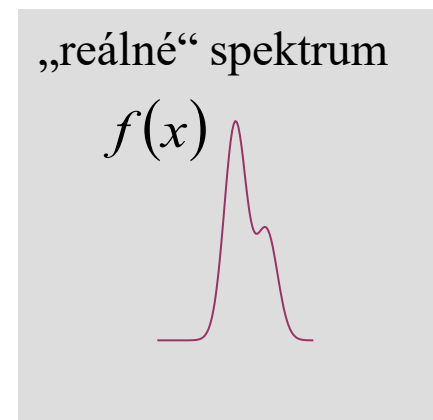
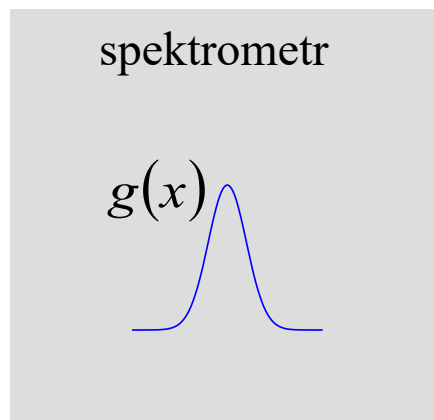
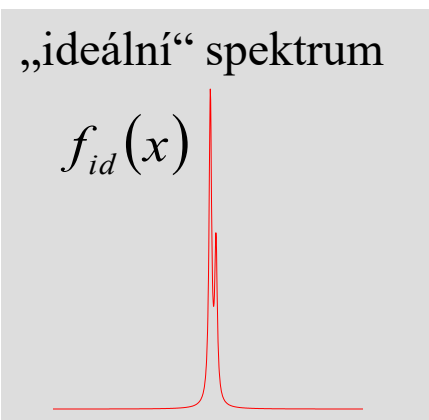
↑  
konvoluce

# Funkce náhodné proměnné – konvoluce

## difrakce rtg. záření



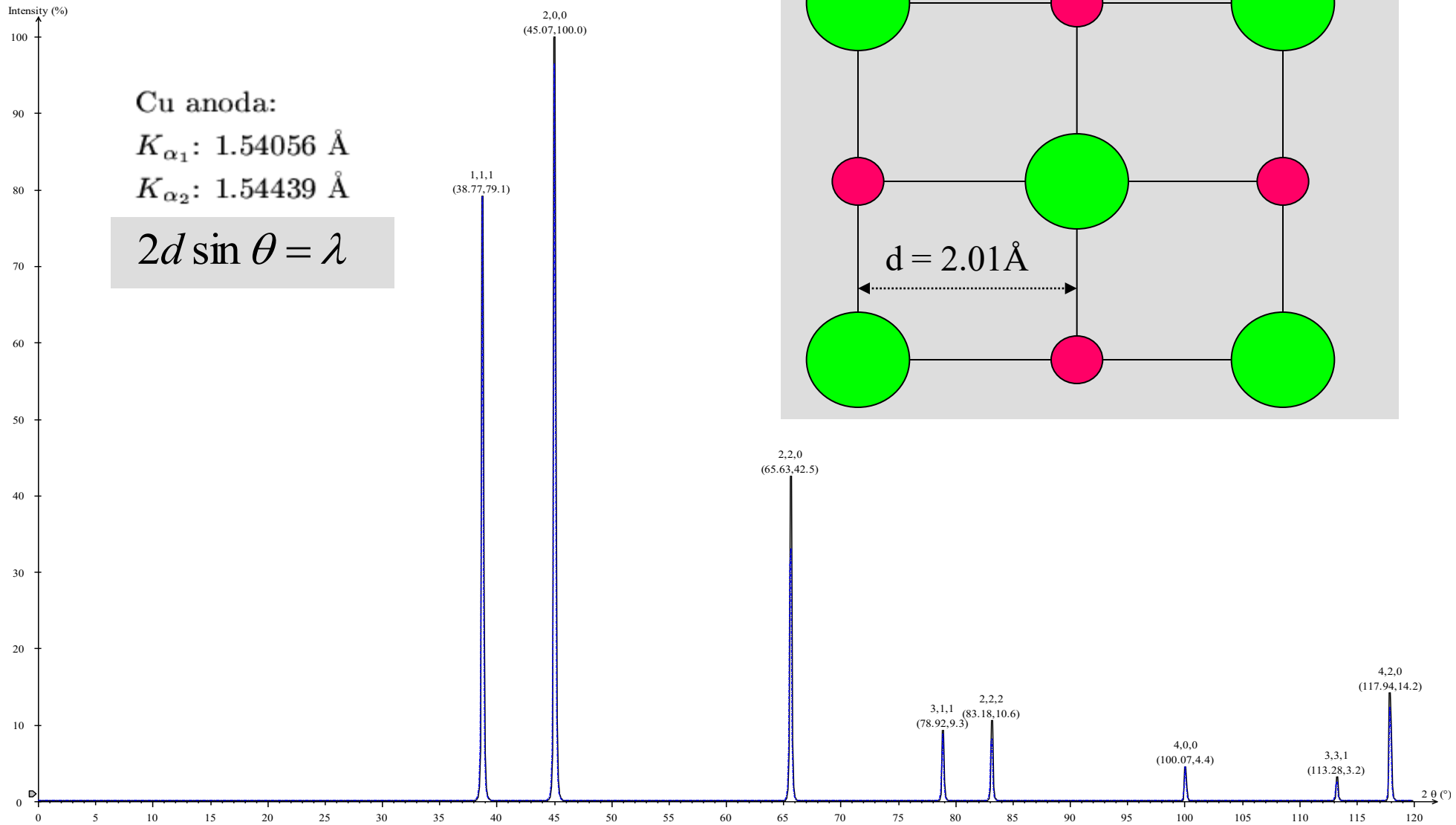
# Funkce náhodné proměnné – konvoluce



$$f(x) = \int_{-\infty}^{\infty} f_{id}(x-y)g(y)dy$$

# Funkce náhodné proměnné – konvoluce

difrakce rtg. záření

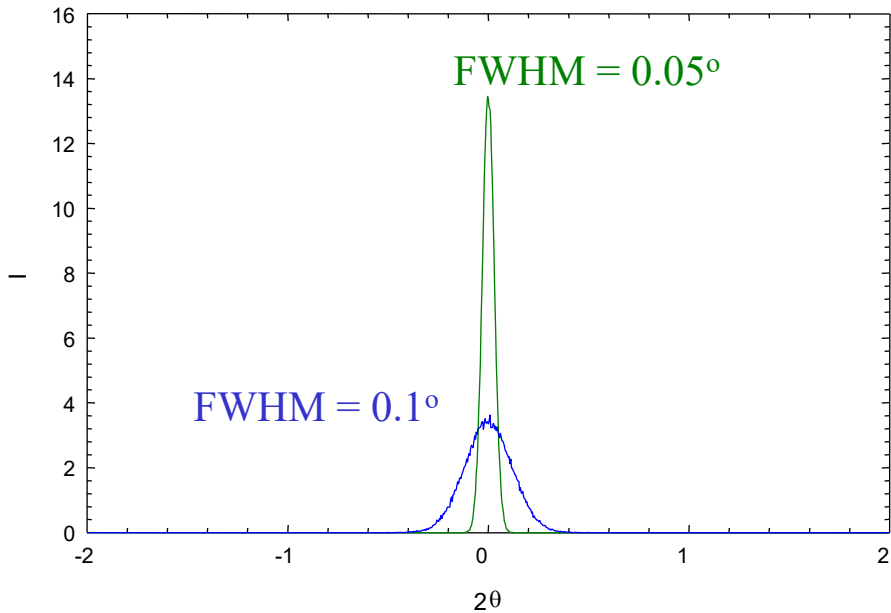


# Funkce náhodné proměnné – konvoluce

Pearson VII

$$I(x) = \frac{I_0}{\left[1 + C(x - x_0)^2\right]^m}$$

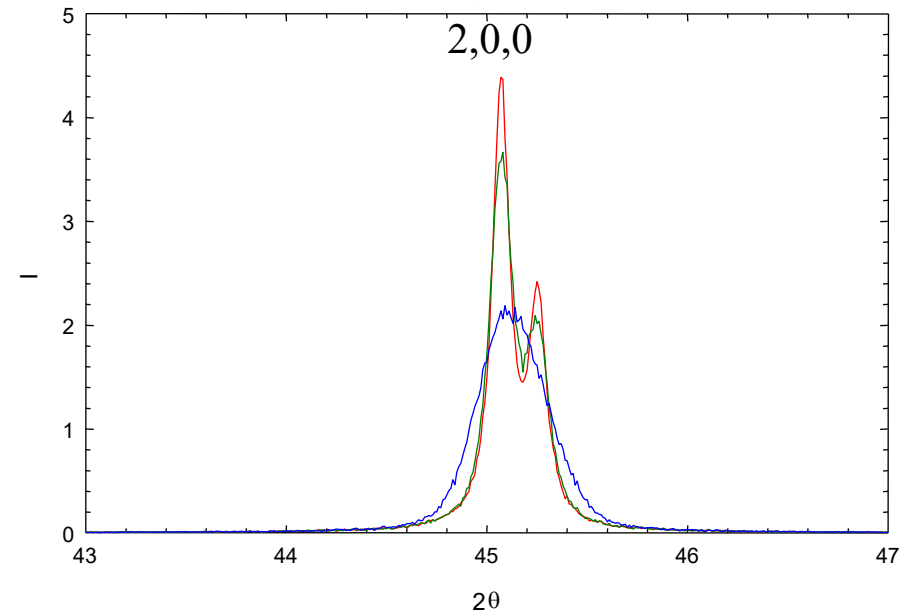
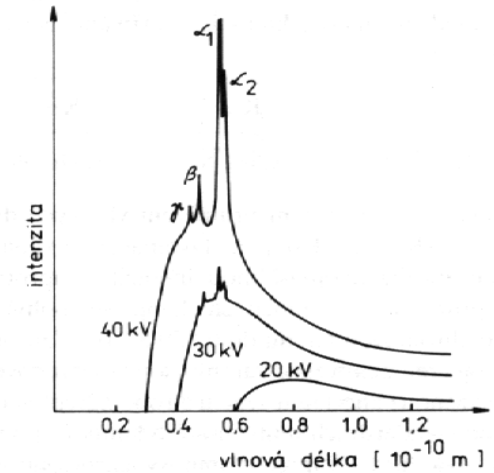
$$C = \frac{\sqrt[m]{2} - 1}{w^2} \quad \text{FWHM} = 2w$$



Cu anoda:

$K_{\alpha_1}$ : 1.54056 Å

$K_{\alpha_2}$ : 1.54439 Å

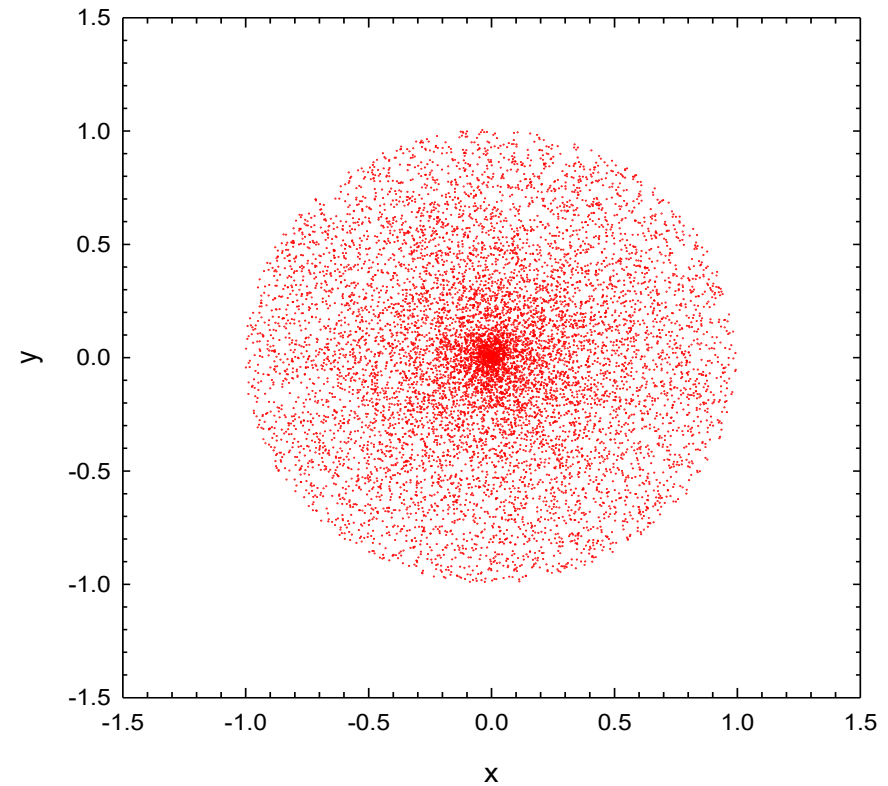
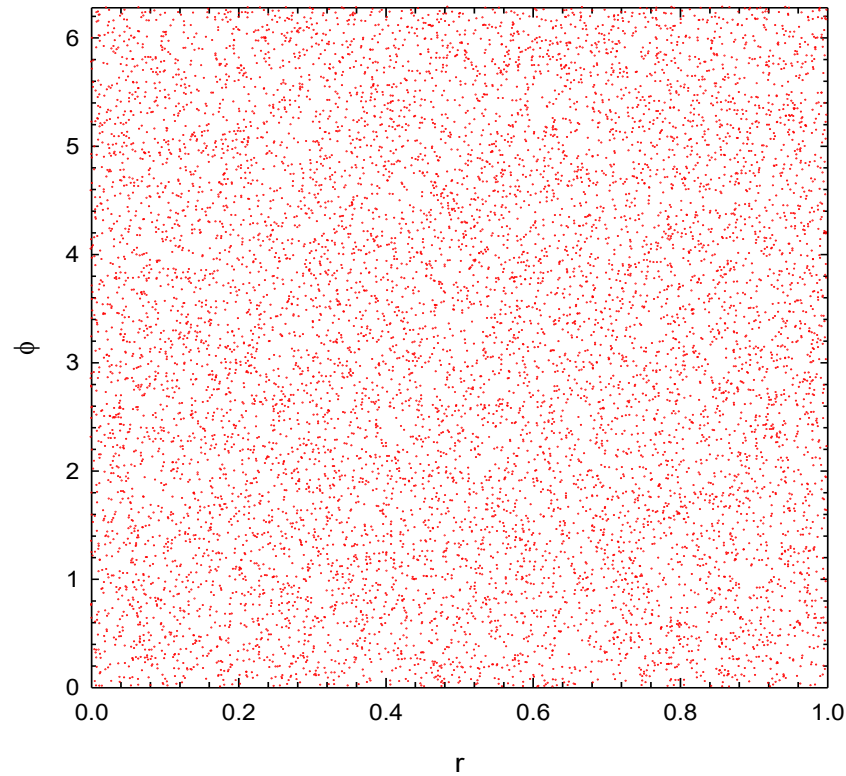


# Funkce náhodné proměnné – transformace souřadnic

$$g(x, y) = \frac{1}{r} f(r, \varphi)$$

$$r \in U(0,1) \quad \varphi \in U(0,2\pi)$$

$$x = r \cos \varphi \quad y = r \sin \varphi$$

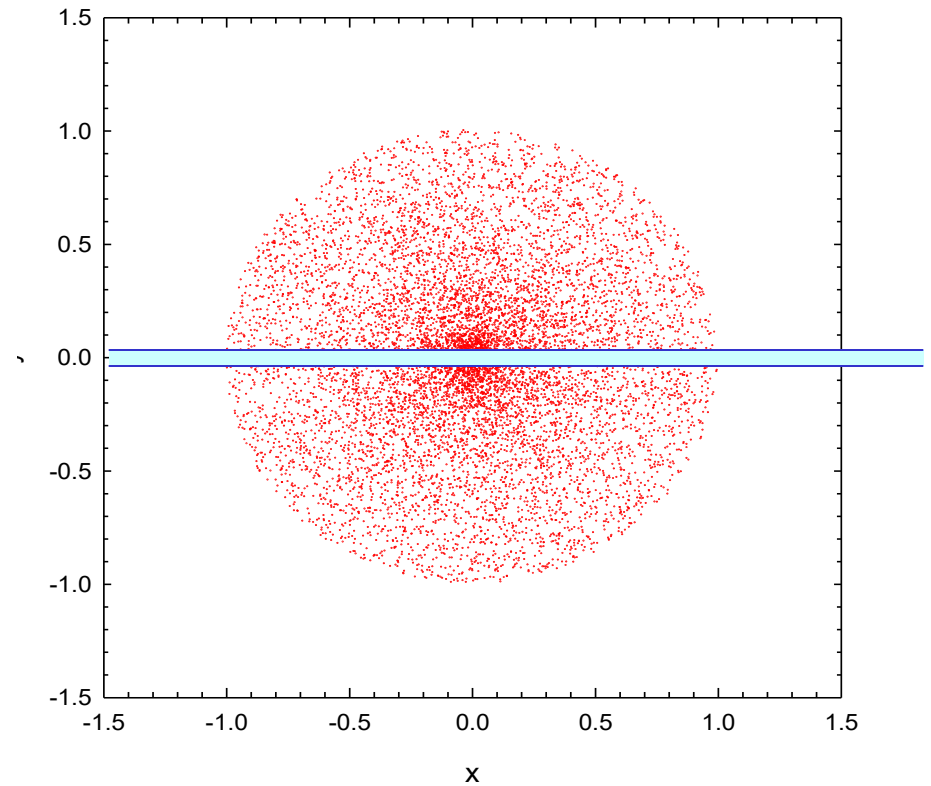
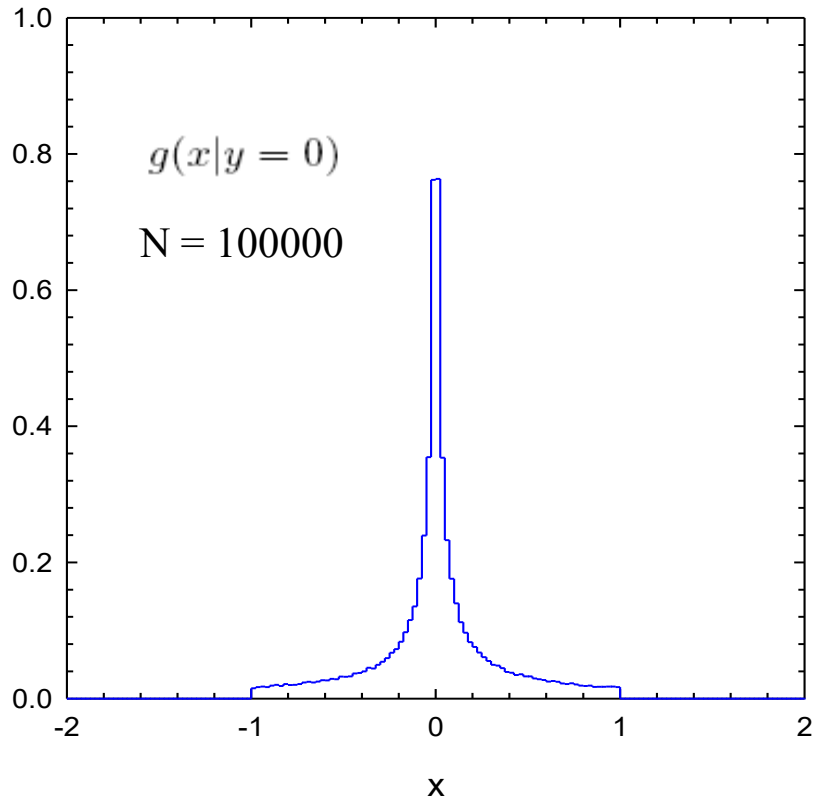


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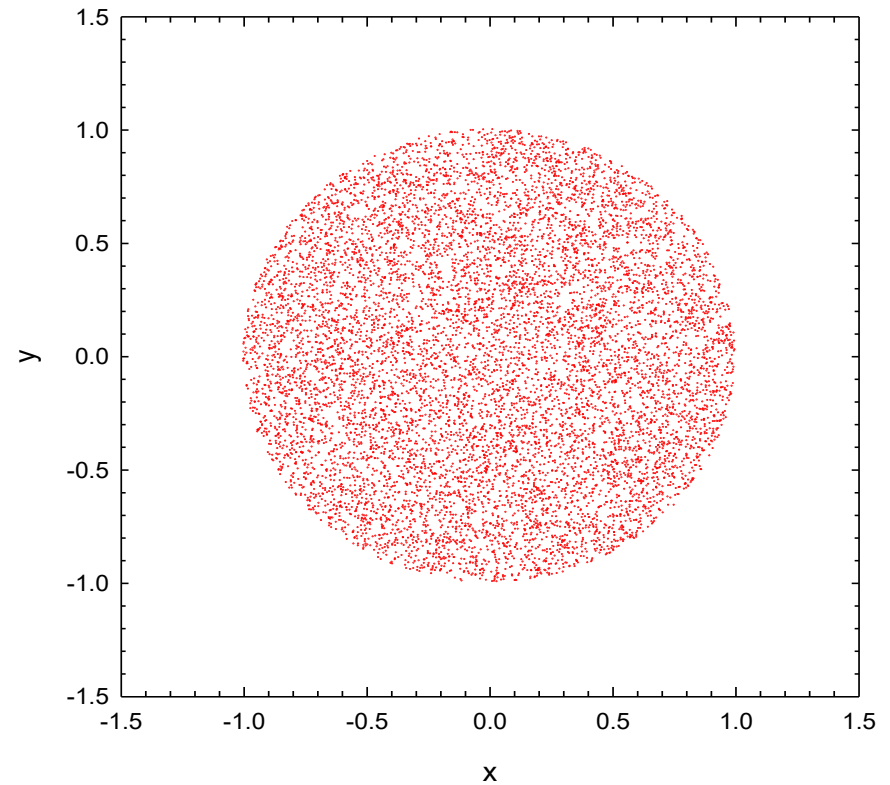
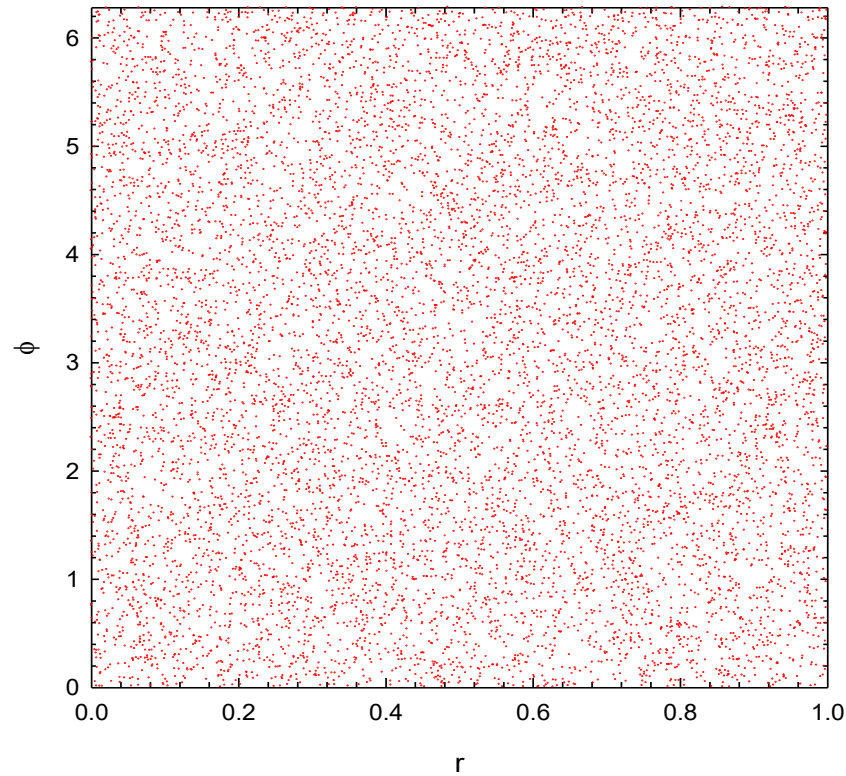


# Funkce náhodné proměnné – transformace souřadnic

$$g(x, y) = 2f(r, \varphi)$$

$$r \in U(0,1) \quad \varphi \in U(0,2\pi)$$

$$x = \sqrt{r} \cos \varphi \quad y = \sqrt{r} \sin \varphi$$

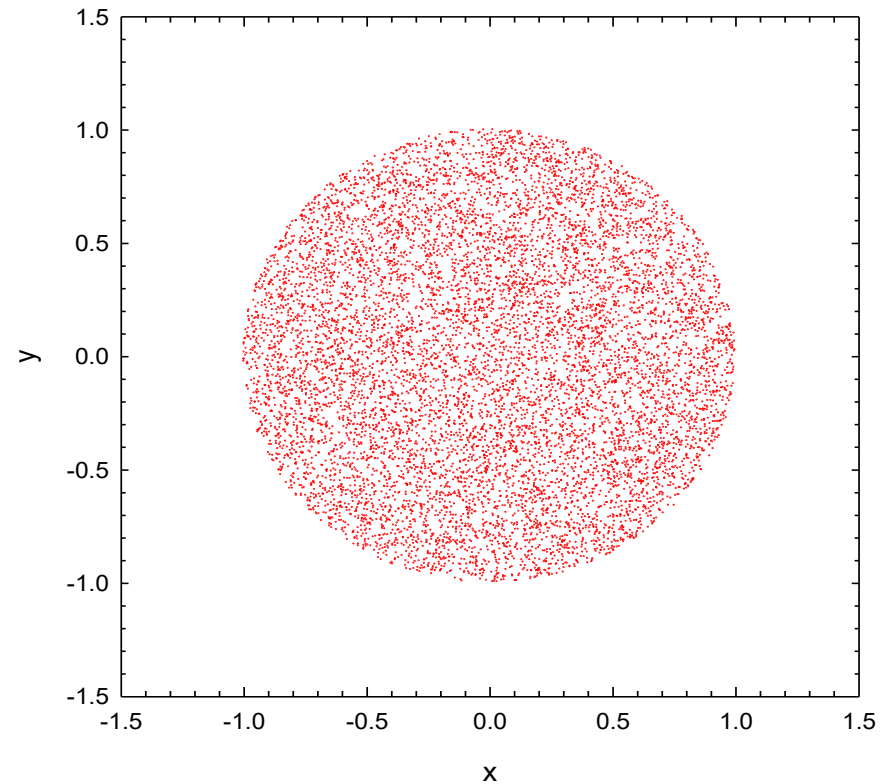
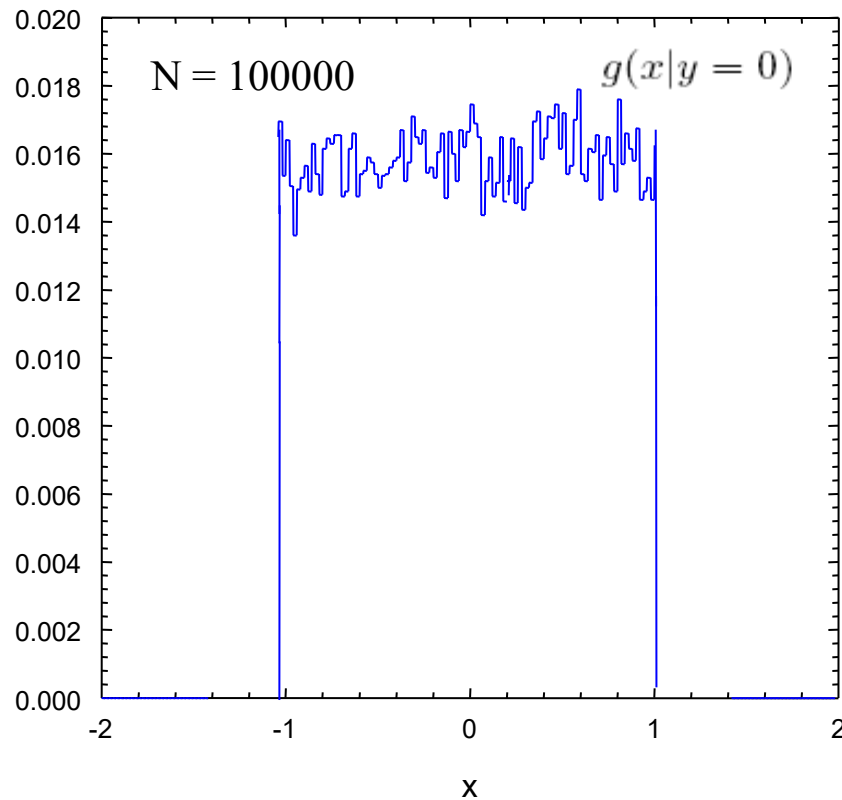


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# Lineární kongruentní generátor

## smíšený generátor

$$I_{j+1} = aI_j + c \pmod{m}$$

$$a, c, m \in \mathbf{N}$$

$$x_j = I_j / m$$

$$0 \leq x_j < 1 \quad x_j \in U(0,1)$$

$$m \approx 2^{32}$$

maximální perioda  $m$

$I_0$  : semínko

korelace

## čistě multiplikativní generátor

$$I_{j+1} = aI_j \pmod{m}$$

$$a, m \in \mathbf{N}$$

$$x_j = I_j / m$$

$$0 \leq x_j < 1 \quad x_j \in U(0,1)$$

$$a = 7^5 = 16807$$

$$m = 2^{31} - 1 = 2147483647$$

$$\text{perioda } 2^{31} - 2 \approx 2.1 \times 10^9$$

# Lineární kongruentní generátor - implementace

## čistě multiplikatívní generátor

$$I_{j+1} = aI_j \pmod{m}$$

$$a, m \in \mathbf{N}$$

$$x_j = I_j / m$$

$$0 \leq x_j < 1 \quad x_j \in U(0,1)$$

$$a = 7^5 = 16807$$

$$m = 2^{31} - 1 = 2147483647$$

$$q = 127773 \quad r = 2836$$

## FaktORIZACE $m$ (Schragerův algoritmus)

$$m = aq + r, \quad q = \lfloor m/a \rfloor, \quad r = m \bmod a$$

$$r < q$$



$$0 < I_j < m - 1 \Rightarrow \begin{cases} 0 \leq a(I_j \bmod q) \leq m - 1 \\ 0 \leq r \lfloor I_j/q \rfloor \leq m - 1 \end{cases}$$

$$aI_j \bmod m = \begin{cases} a(I_j \bmod q) - r \lfloor I_j/q \rfloor, & \text{je-li tento výraz} \geq 0 \\ a(I_j \bmod q) - r \lfloor I_j/q \rfloor + m, & \text{jinak} \end{cases}$$

# Lineární kongruentní generátor - implementace

```
float ran0(int *p_i0)
{
    #define a (16807)
    #define m (2147483647)
    #define q (127773)
    #define r (2836)

    int k,i0;
    float x;

    i0=*p_i0;
    k=i0/q; // [I0/q]
    i0=a*(i0-k*q)-r*k; // a(I0 mod q)-r[I0/q]
    if(i0<0) i0=i0+m;
    x=(float)i0/m; //converze na realne cislo z intervalu (0,1)
    *p_i0=i0;
    return(x);
}
```

# Lineární kongruentní generátor - implementace

```
void main()
{
    FILE *f,*g;
    int iseed;
    int i;
    float x;

    if((f=fopen("iseed.dat","r"))==NULL) iseed=123456789;
        else
        {
            f=fopen("iseed.dat","r");
            fscanf(f,"%d",&iseed);
            fclose(f);
        }

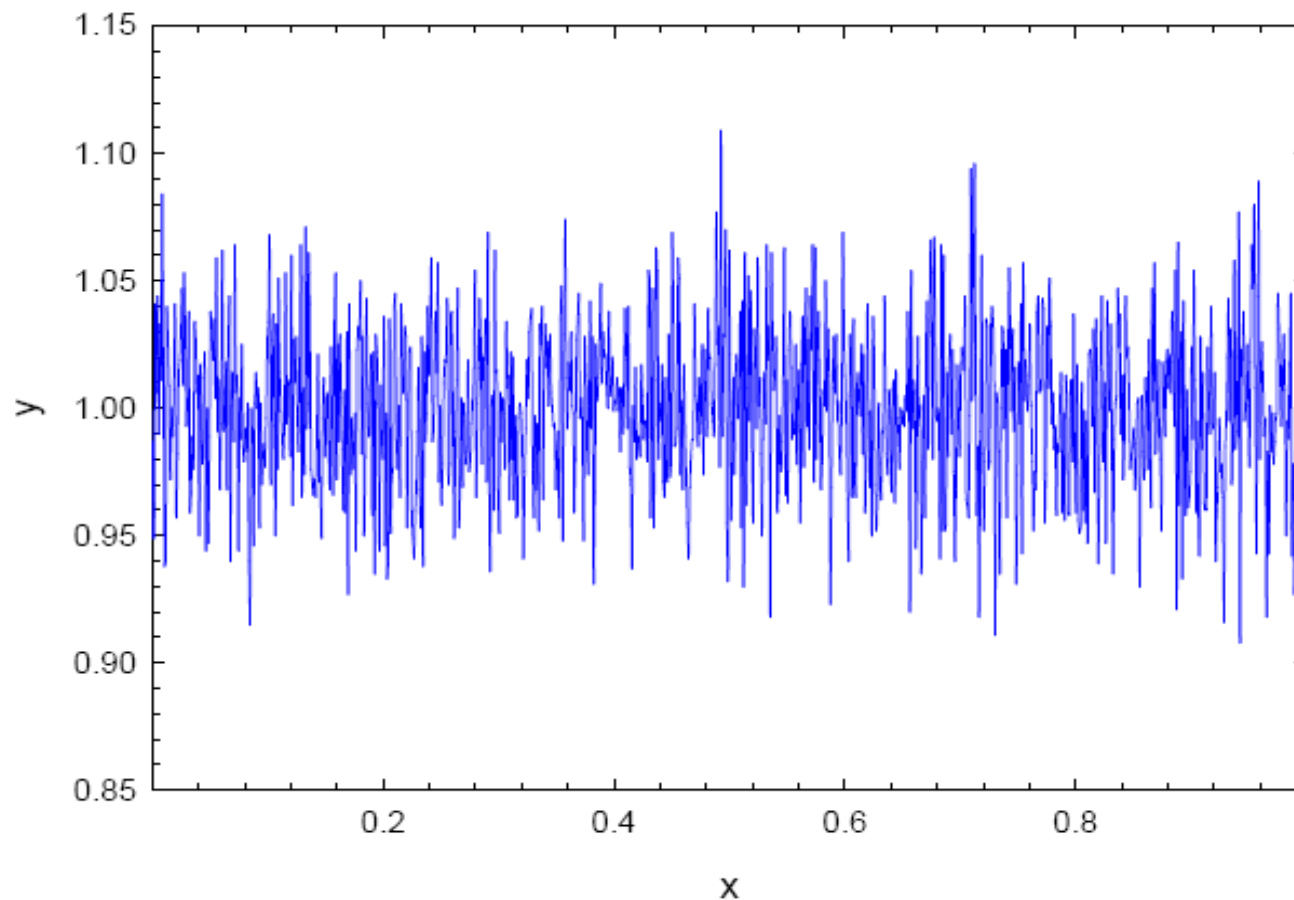
    g=fopen("ran0.dat","w");
    for(i=0;i<1000000;i++)
    {
        x=ran0(&iseed);
        fprintf(g,"%f\n",x);
    }

    f=fopen("iseed.dat","w");
    fprintf(f,"%d",iseed);
    fclose(f);
    fclose(g);
}
```

# Lineární kongruentní generátor - implementace

$N = 1000000$

multiplikativní generátor,  $a = 16807$ ,  $m = 2^{31} - 1$



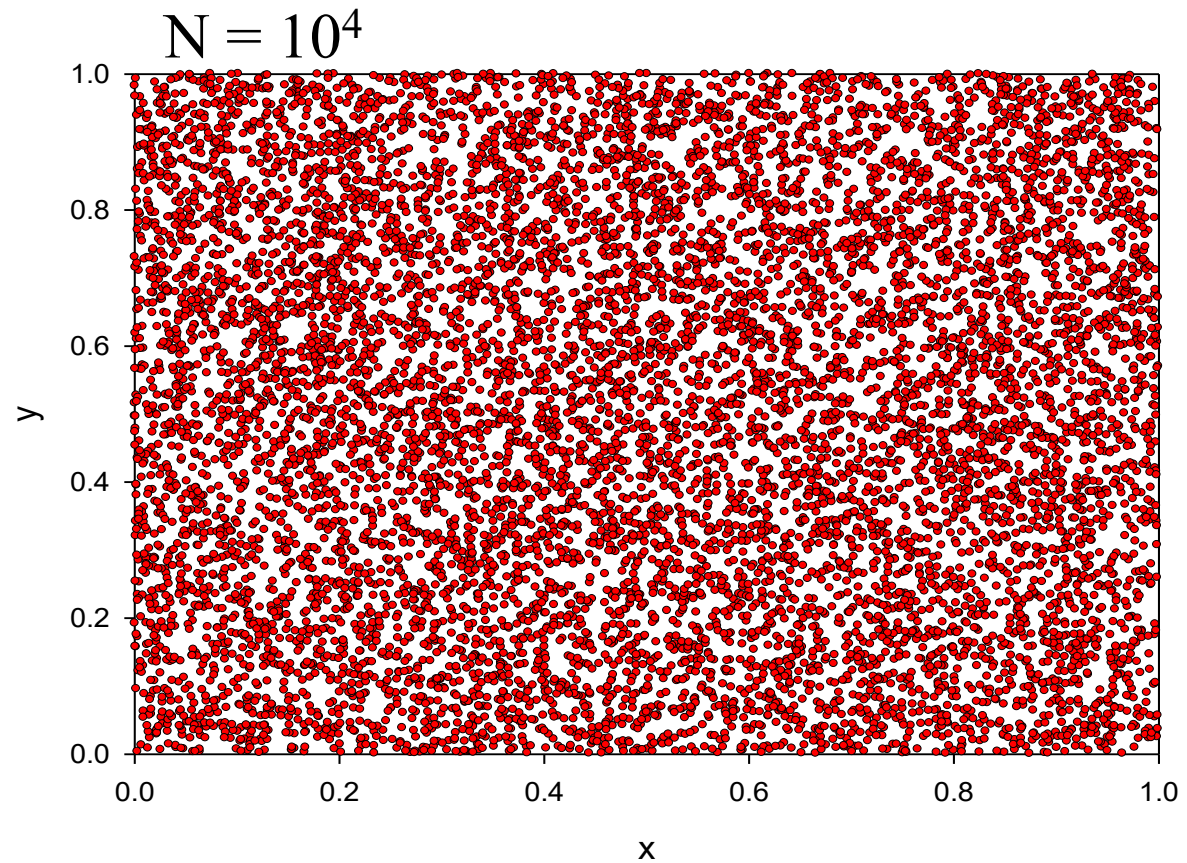
# Lineární kongruentní generátor – test

IBM RANDU

$$I_{j+1} = a I_j \pmod{m}$$

$$a = 65539$$

$$m = 2^{31}$$





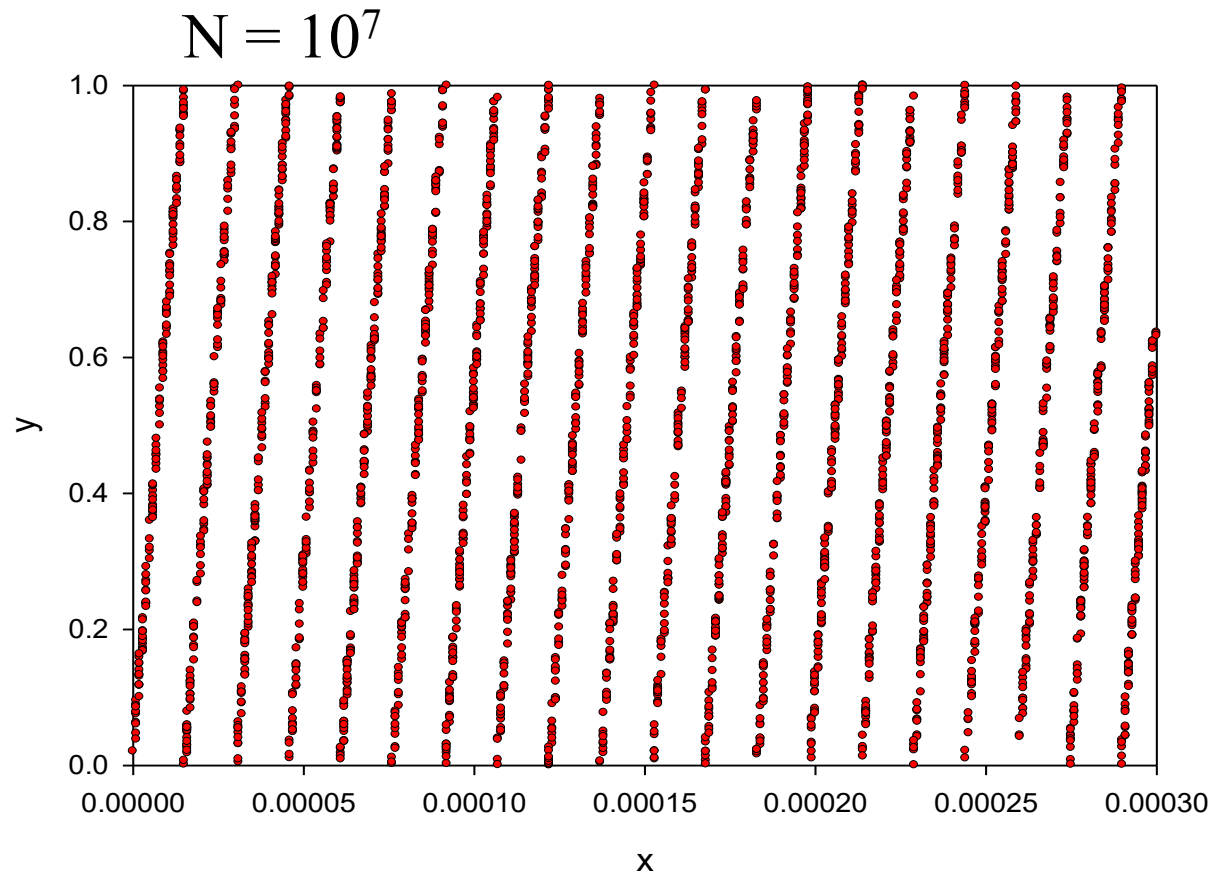
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# Lineární kongruentní generátor – sériová korelace

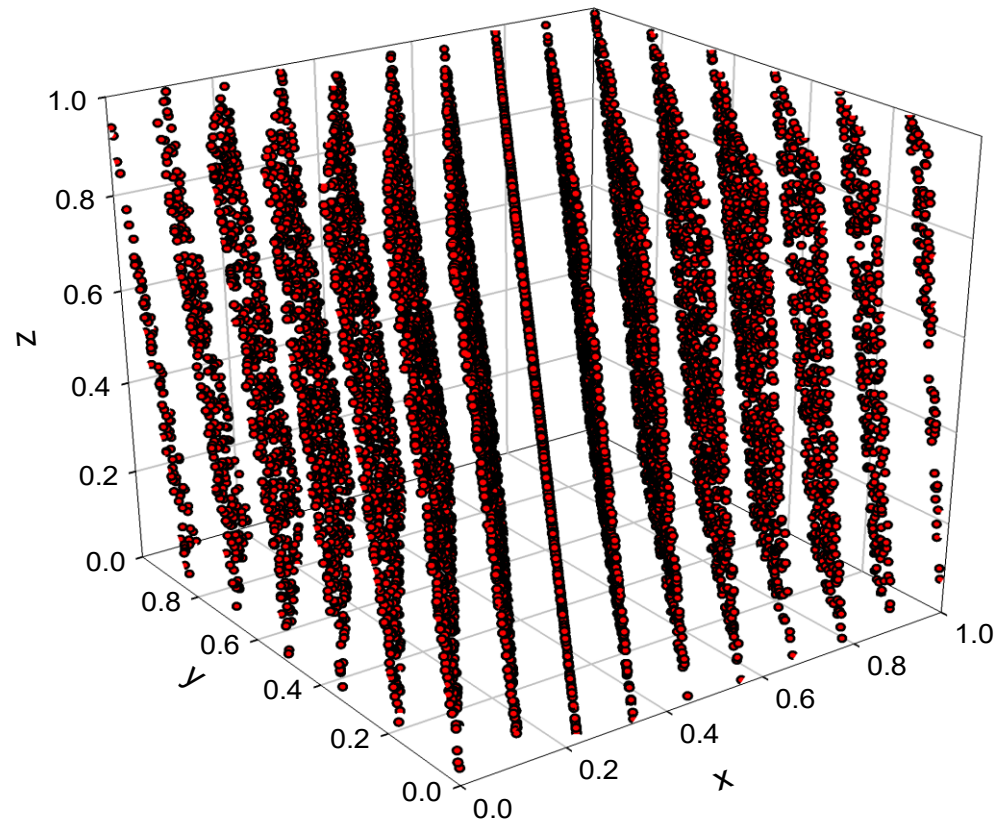
IBM RANDU

$$I_{j+1} = a I_j \pmod{m}$$

$$N = 10^4$$

$$a = 65539$$

$$m = 2^{31}$$



“ We guarantee that each number is random individually,  
but we don't guarantee that more than one of them is random.”

# Lineární kongruentní generátor – sériová korelace

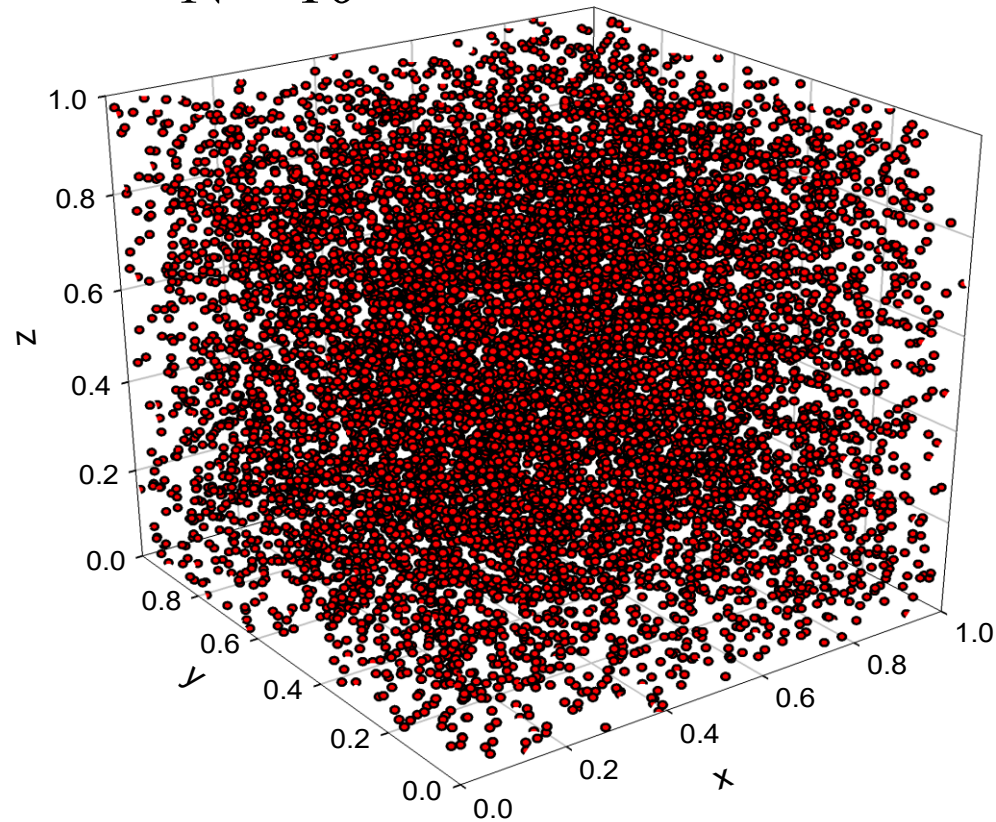
RAN0

$$I_{j+1} = a I_j \pmod{m}$$

$$a = 16807$$

$$m = 2^{31} - 1$$

$N = 10^4$



# Lineární kongruentní generátor – posuvný registr

RAN2 – L'Ecuyer

$$I_{j+1} = a I_j \pmod{m}$$

$$a_1 = 40014$$

$$m_1 = 2147483563$$

$$a_2 = 40692$$

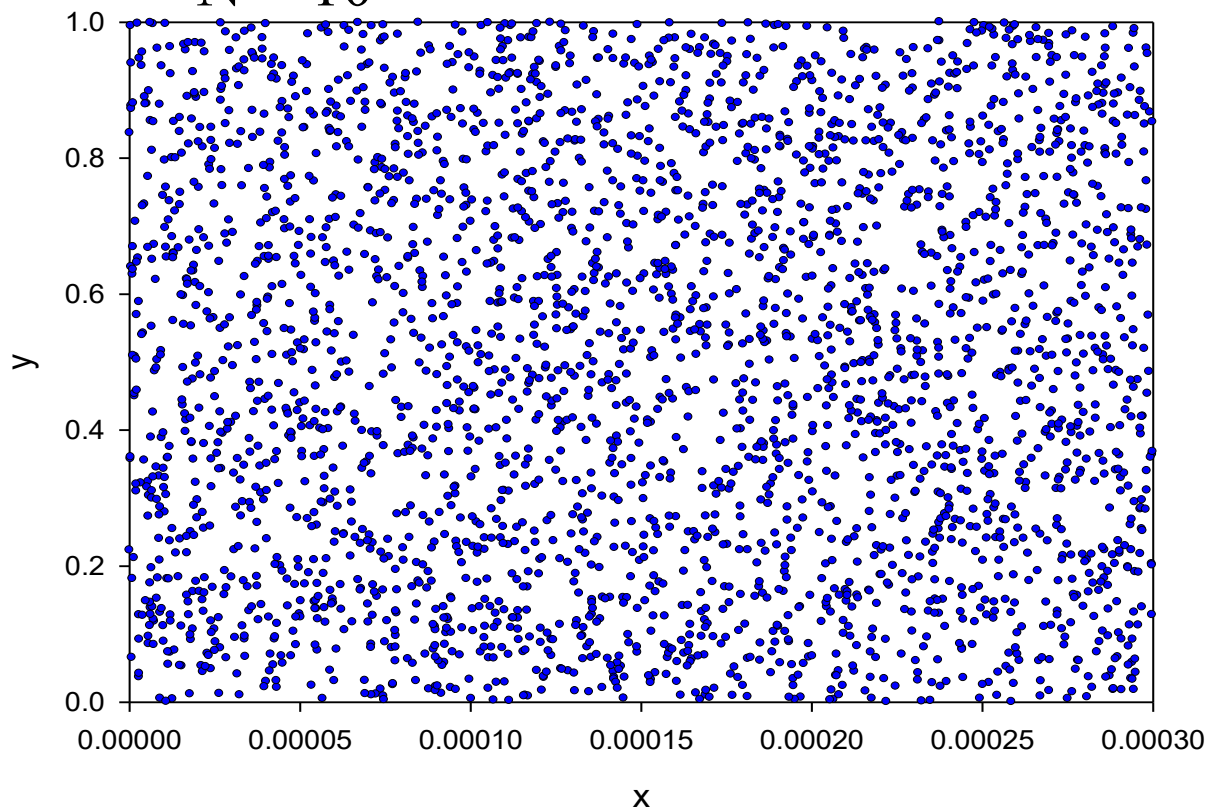
$$m_2 = 2147483399$$

+ posuv registru

perioda  $\approx 2.3 \times 10^{18}$

$$I_j = I_{j,1} + I_{j,2} \pmod{m_1 + m_2}$$

$N = 10^7$



# Lineární kongruentní generátor – barevný test

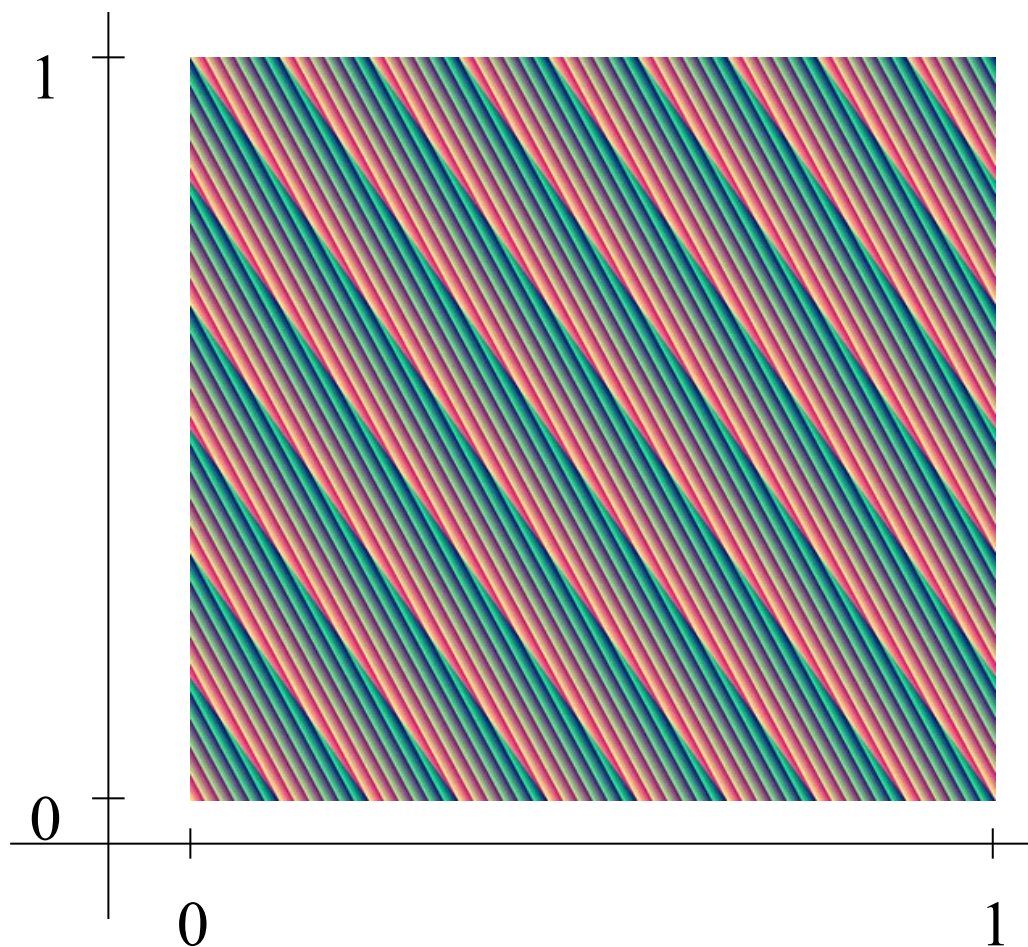
IBM RANDU

$$I_{j+1} = a I_j \pmod{m}$$

$$a = 65539$$

$$m = 2^{31}$$

ze dvou čísel generován bod ve čtverci  $[0,1] \times [0,1]$   
barva náhodně ze stejného generátoru, například  
pomocí RGB složek v rozsahu 0-255





# Lineární kongruentní generátor – barevný test

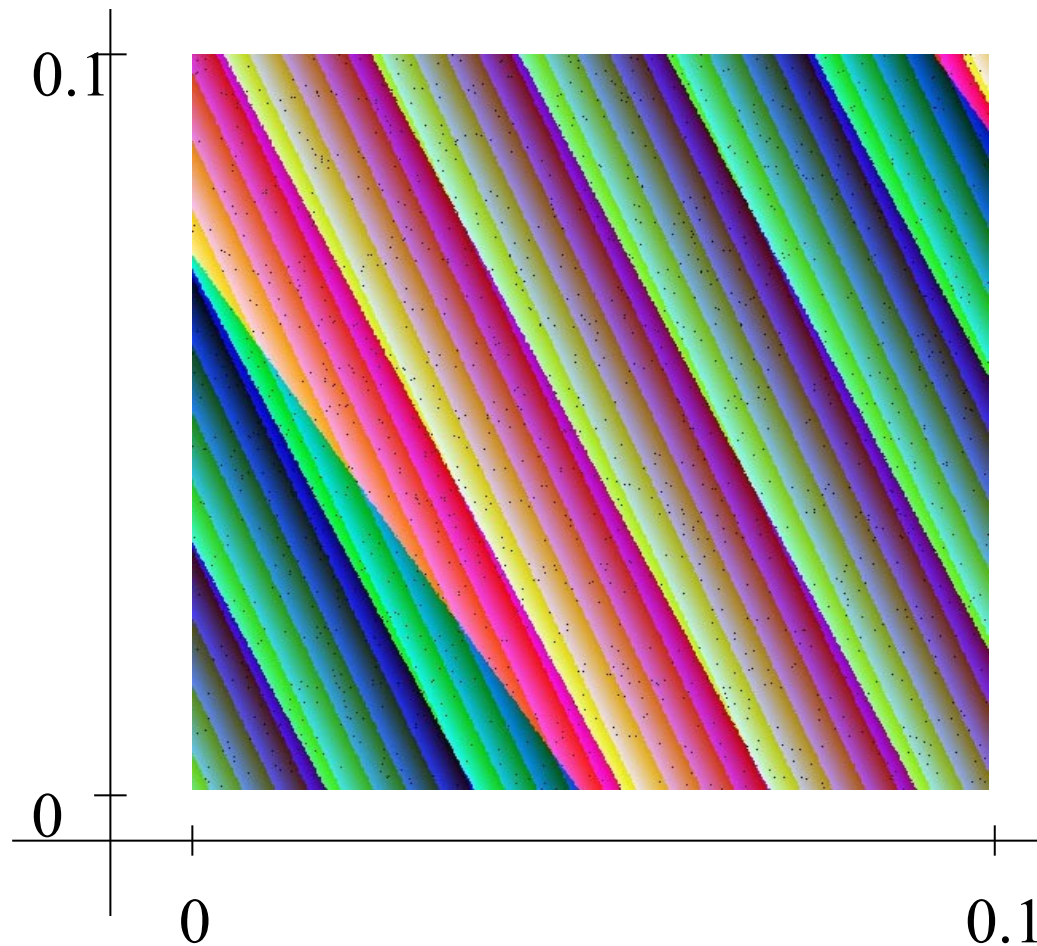
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# Lineární kongruentní generátor – barevný test

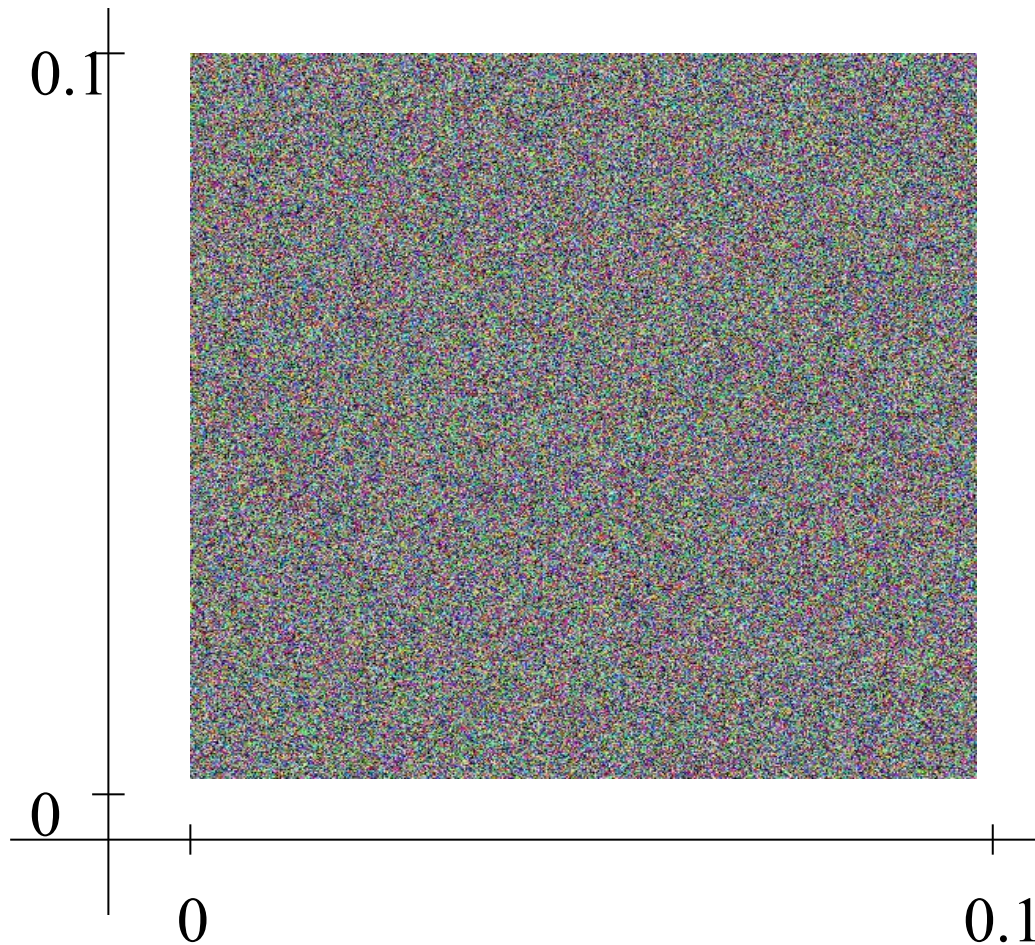
RAND()

Mersenne twister

perioda  $2^{19937}-1$

MS Visual C++ 6.0

ze dvou čísel generován bod ve čtverci  $[0,1] \times [0,1]$   
barva náhodně ze stejného generátoru, například  
pomocí RGB složek v rozsahu 0-255



# Generátory pseudonáhodných čísel

	Statistical Quality	Prediction Difficulty	Reproducible Results	Multiple Streams	Period	Useful Features	Time Performance	Space Usage	Code Size & Complexity	$k$ -Dimensional Equidistribution
<b>PCG Family</b>	Excellent	Challenging	Yes	Yes (e.g. $2^{63}$ )	Arbitrary	Jump ahead, Distance	Very fast	Very compact	Very small	Arbitrary*
<b>Mersenne Twister</b>	Some Failures	Easy	Yes	No	Huge $2^{19937-1}$	Jump ahead	Acceptable	Huge (2 KB)	Complex	623
<b>Arc4Random</b>	Some Issues	Secure	Not Always	No	Huge $2^{1699}$	No	Slow	Large (0.5 KB)	Complex	No
<b>ChaCha20<sup>†</sup></b>	Good	Secure	Yes	Yes ( $2^{128}$ )	$2^{128}$	Jump ahead, Distance	Fairly Slow	Plump (0.1 KB)	Complex	No
<b>Minstd (LCG)</b>	Many Issues	Trivial	Yes	No	Tiny $< 2^{32}$	Jump ahead, Distance	Acceptable	Very compact	Very small	No
<b>LCG 64/32</b>	Many Issues	Published Algorithms	Yes	Yes $2^{63}$	Okay $2^{64}$	Jump ahead, Distance	Very fast	Very compact	Very small	No
<b>XorShift 32</b>	Many Issues	Trivial	Yes	No	Small $2^{32}$	Jump ahead	Fast	Very compact	Very small	No
<b>XorShift 64</b>	Many Issues	Trivial	Yes	No	Okay $2^{64}$	Jump ahead	Fast	Very compact	Very small	No
<b>RanQ</b>	Some Issues	Trivial	Yes	No	Okay $2^{64}$	Jump ahead	Fast	Very compact	Very small	No
<b>XorShift* 64/32</b>	Excellent	Unknown?	Yes	No	Okay $2^{64}$	Jump ahead	Fast	Very compact	Very small	No



# Lineární kongruentní generátor – barevný test

PCG-64 (Permutation congruential generator)

Perioda  $2^{138}$

Python

	Statistical Quality	Prediction Difficulty	Reproducible Results	Multiple Streams	Period	Useful Features	Time Performance	Space Usage	Code Size & Complexity	$k$ -Dimensional Equidistribution
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